

First-Order Theorem Proving

Vampire Cookies

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From theory to practice

- ▶ Preprocessing and CNF transformation;
- ▶ Superposition system;
- ▶ Orderings;
- ▶ Selection functions;
- ▶ Fairness (saturation algorithms);
- ▶ Redundancy.

Vampire's preprocessing (incomplete list)

1. (Optional) Select a **relevant subset** of formulas.
2. (Optional) Add **theory axioms**;
3. **Rectify** the formula.
4. If the formula contains any occurrence of \top or \perp , **simplify** the formula.
5. Remove **if-then-else** and **let-in** connectives.
6. Apply **pure predicate elimination**.
7. (Optional) Remove **unused predicate definitions**.
8. Convert the formula into **equivalence negation normal form (ENNF)**.
9. Use a **naming technique** to replace some subformulas by their names.
10. Convert the formula into **negation normal form (NNF)**.
11. **Skolemize** the formula.
12. (Optional) Replace **equality axioms**.
13. Determine a **literal ordering** to be used.
14. Transform the formula into its **clausal normal form**.
15. Remove **tautologies**.
16. **Pure literal elimination**.

How to Design a Good Saturation Algorithm?

A saturation algorithm must be **fair**: every possible generating inference must eventually be selected.

Two main implementation principles:

apply simplifying inferences
eagerly;
apply generating inferences
lazily.

checking for simplifying
inferences should pay off;
so it must be cheap.

Given Clause Algorithm (no Simplification)

```
input: init: set of clauses;  
var active, passive, queue: sets of clauses;  
var current: clauses ;  
active :=  $\emptyset$ ;  
passive := init;  
while passive  $\neq \emptyset$  do  
*   current := select(passive);  
    move current from passive to active;  
*   queue := infer(current, active);  
    if  $\square \in$  queue then return unsatisfiable;  
    passive := passive  $\cup$  queue  
od;  
return satisfiable
```

(* clause selection *)

(* generating inferences *)

Given Clause Algorithm (with Simplification)

In fact, there is more than one ...

Otter Saturation Algorithm

input: *init*: set of clauses;

var *active*, *passive*, *unprocessed*: set of clauses;

var *given*, *new*: clause;

active := \emptyset ;

unprocessed := *init*;

loop

while *unprocessed* $\neq \emptyset$

new := *pop*(*unprocessed*);

if *new* = \square **then return** *unsatisfiable*;

if *retained*(*new*) **then**

 (* retention test *)

 simplify *new* by clauses in *active* \cup *passive*; (* forward simplification *)

if *new* = \square **then return** *unsatisfiable*;

if *retained*(*new*) **then**

 (* another retention test *)

 delete and simplify clauses in *active* and (* backward simplification *)

passive using *new*;

 move the simplified clauses to *unprocessed*;

 add *new* to *passive*

if *passive* = \emptyset **then return** *satisfiable* or *unknown*

given := *select*(*passive*);

 (* clause selection *)

 move *given* from *passive* to *active*;

unprocessed := *infer*(*given*, *active*);

 (* generating inferences *)

Age-Weight Ratio

How to select nice clauses?

- ▶ Small clauses are nice.
- ▶ Selecting only small clauses can postpone the selection of an old clause (e.g., input clause) for too long, in practice resulting in incompleteness.

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Solution:

- ▶ A fixed percentage of clauses is selected **by weight**, the rest are selected **by age**.
- ▶ So we use an **age-weight ratio** $a : w$: of each $a + w$ clauses select a **oldest** and w **smallest** clauses.

Limited Resource Strategy

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Try:

```
vampire --age_weight_ratio 4:1
  --forward_subsumption_resolution off
  --time_limit 20
  GRP140-1.p
```

CASC Mode

```
vampire --mode casc SET014-3.p
```